



UNIVERSITI PUTRA MALAYSIA

***PERMISSION-BASED FAULT TOLERANT MUTUAL EXCLUSION
ALGORITHM FOR MOBILE AD HOC NETWORKS***

FARANEH ZARAFSHAN

FK 2015 89



**PERMISSION-BASED FAULT TOLERANT MUTUAL EXCLUSION
ALGORITHM FOR MOBILE AD HOC NETWORKS**

By

FARANEH ZARAFSHAN

**Thesis Submitted to the School of Graduate Studies, Universiti Putra Malaysia,
in Fulfilment of the Requirements for the Degree of Doctor of Philosophy**

May 2015

All material contained within the thesis, including without limitation text, logos, icons, photographs and all other artwork, is copyright material of Universiti Putra Malaysia unless otherwise stated. Use may be made of any material contained within the thesis for non-commercial purposes from the copyright holder. Commercial use of material may only be made with the express, prior, written permission of Universiti Putra Malaysia.

Copyright © Universiti Putra Malaysia



This thesis is dedicated to

My husband Abbas

For teaching me not to be disappointed and taking steps with me in hard moments of
life

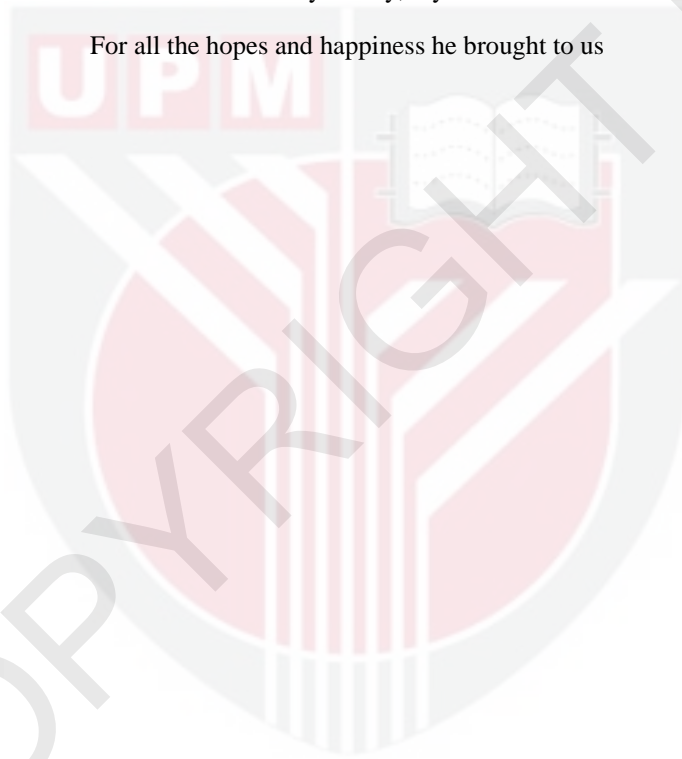
My dear parents

For their endless support and love

and

New member of my family, my son: Radwin

For all the hopes and happiness he brought to us



Abstract of thesis presented to the Senate of Universiti Putra Malaysia in fulfilment of the requirement for the degree of Doctor of Philosophy

PERMISSION-BASED FAULT TOLERANT MUTUAL EXCLUSION ALGORITHM FOR MOBILE AD HOC NETWORKS

By

FARANEH ZARAFSHAN

May 2015

Chairman: Syed Abdul Rahman Al-Haddad B. Syed Mohamed, PhD
Faculty: Engineering

This study focuses on resolving the problem of mutual exclusion in mobile ad hoc networks. A Mobile Ad Hoc Network (MANET) is a wireless network without fixed infrastructure. Nodes are mobile and topology of MANET changes very frequently and unpredictably. Due to these limitations, conventional mutual exclusion algorithms presented for distributed systems (DS) are not applicable for MANETs unless they attach to a mechanism for dynamic changes in their topology.

Algorithms for mutual exclusion in DS are categorized into two main classes including token-based and permission-based algorithms. Token-based algorithms depend on circulation of a specific message known as token. The owner of the token has priority for entering the critical section. Token may lose during communications, because of link failure or failure of token host. However, the processes for token-loss detection and token regeneration are very complicated and time-consuming. Token-based algorithms are generally non-fault-tolerant (although some mechanisms are utilized to increase their level of fault-tolerance) because of common problem of single token as a single point of failure. On the contrary, permission-based algorithms utilize the permission of multiple nodes to guarantee mutual exclusion. It yields to high traffic when number of nodes is high. Moreover, the number of message transmissions and energy consumption increase in MANET by increasing the number of mobile nodes accompanied in every decision making cycle.

The purpose of this study is to introduce a method of managing the critical section, named as Ancestral, having higher fault-tolerance than token-based and fewer message transmissions and traffic rather than permission-based algorithms. This method makes a tradeoff between token-based and permission-based. It does not utilize any token, that is similar to permission-based, and the latest node having the critical section influences the entrance of the next node to the critical section, that is similar to token-based algorithms. The algorithm based on ancestral is named as DAD algorithms and increases the availability of fully connected network between 2.86 to 59.83% and

decreases the number of message transmissions from $4j-2$ to $3j$ messages (j as number of nodes in partition).

This method is then utilized as the basis of dynamic ancestral mutual exclusion algorithm for MANET which is named as MDA. This algorithm is presented and evaluated for different scenarios of mobility of nodes, failure, load and number of nodes. The results of study show that MDA algorithm guarantees mutual exclusion, dead lock freedom and starvation freedom. It improves the availability of CS to minimum 154.94% and 113.36% for low load and high load of CS requests respectively compared to other permission-based algorithm. Furthermore, it improves response time up to 90.69% for high load and 75.21% for low load of CS requests. It degrades the number of messages from n to 2 messages in the best case and from $3n/2$ to n in the worst case. MDA algorithm is resilient to transient partitioning of network that is normally occurs due to failure of nodes or links.

Abstrak tesis yang dikemukakan kepada Senat Universiti Putra Malaysia sebagai memenuhi keperluan untuk ijazah Doktor Falsafah

ALGORITMA SALING PENGASINGAN BERASASKAN PERMINTAAN TOLERANSI RALAT UNTUK RANGKAIAN AD HOC MUDAH ALIH

Oleh

FARANEH ZARAFSHAN

Mei 2015

Pengerusi: Syed Abdul Rahman Al-Haddad B. Syed Mohamed, PhD
Fakulti: Kejuruteraan

Tesis ini memfokuskan kepada menyelesaikan masalah saling pengasingan dalam rangkaian talian sementara mudah alih. Rangkaian Talian Sementara Mudah Alih (MANET) adalah sebuah rangkaian wayarles tanpa infrastruktur yang tetap. Nod-nod adalah mudah alih. Topologi MANET sangat kerap berubah dan tidak menentu. Disebabkan batasan-batasan ini, algoritma saling pengasingan yang biasanya dikemukakan untuk sistem teragih (DS) tidak boleh digunakan untuk MANETs melainkan mereka disambungkan kepada satu mekanisme untuk perubahan dinamik di dalam topologi mereka.

Algoritma untuk saling pengasingan dalam DS dikategorikan kepada dua kelas utama yang merangkumi algoritma berasaskan-token dan berasaskan-kebenaran. Algoritma berasaskan-token bergantung kepada peredaran sebuah mesej khusus yang dikenali sebagai token. Pemilik token itu mempunyai keutamaan untuk memasuki bahagian yang kritikal. Token boleh hilang sewaktu komunikasi, disebabkan kegagalan sambungan atau kegagalan hos token. Walau bagaimanapun, proses-proses untuk pengesanan token-yang-hilang dan penjanaan semula token adalah sangat rumit dan memakan masa. Algoritma berasaskan-token biasanya bukan tahan-rosak (walaupun beberapa mekanisme telah digunakan untuk meningkatkan tahap tahan-kerosakan mereka) disebabkan oleh masalah biasa token tunggal sebagai kegagalan titik tunggal. Sebaliknya, algoritma berasaskan-kebenaran menggunakan kebenaran beberapa nod untuk menjamin saling pengasingan. Oleh itu, tiada kegagalan token dan tiada kegagalan titik tunggal. Walau bagaimanapun, ia menyumbang kepada trafik yang tinggi apabila bilangan nod bertambah. Tambahan pula, bilangan penghantaran mesej dan penggunaan tenaga dalam MANET meningkat dengan bertambahnya bilangan nod mudah alih yang mengiringi setiap kitaran membuat keputusan.

Tujuan tesis ini adalah untuk memperkenalkan satu kaedah baru dalam menguruskan bahagian kritikal, yang dinamakan sebagai Warisan, yang mempunyai tahan-kerosakan yang lebih tinggi daripada berasaskan-token dan mempunyai penghantaran mesej dan trafik yang lebih rendah berbanding algoritma berasaskan-kebenaran.

Kaedah ini memberi keseimbangan antara berasaskan-token dan berasaskan-kebenaran. Ia tidak menggunakan apa-apa token, iaitu serupa dengan berasaskan-kebenaran, dan nod terbaru yang mempunyai bahagian yang kritikal mempengaruhi kemasukan nod seterusnya ke bahagian kritikal, iaitu serupa dengan algoritma berasaskan-token. Algoritma yang berasaskan warisan dinamakan sebagai algoritma DAD dan meningkatkan ketersediaan rangkaian bersambung penuh antara 2.86 sehingga 59.83% dan mengurangkan bilangan penghantaran mesej dari $4j-2$ sehingga $3j$ mesej (j sebagai bilangan nod-nod didalam pemetakan).

Kaedah baru kemudiannya digunakan sebagai asas kepada algoritma saling pengasingan warisan dinamik untuk MANET. Algoritma ini dibentangkan dan dinilai untuk senario-senario yang berbeza dari segi mobiliti nod, kegagalan, beban dan bilangan nod. Keputusan kajian menunjukkan bahawa algoritma MDA menjamin saling pengasingan, kebebasan kunci mati dan kebebasan kebuluran. Ia juga memperbaiki ketersediaan CS kepada minimum 154.94% dan 113.36% masing-masing untuk permintaan-permintaan CS bebanan rendah dan bebanan tinggi berbanding dengan algoritma berasaskan-kebenaran yang lain. Tambahan pula, ia memperbaiki masa tindak balas sehingga 90.69% untuk bebanan tinggi dan 75.21% untuk permintaan-permintaan CS bebanan rendah. Ia mengurangkan bilangan mesej dari n kepada 2 mesej dalam kes yang terbaik dan dari $3n/2$ kepada n dalam kes yang paling teruk. Algoritma MDA mempunyai ketahanan tinggi kepada pemetakan sementara rangkaian-rangkaian yang biasanya berlaku disebabkan oleh kegagalan nod atau sambungan.

ACKNOWLEDGEMENTS

I would like to thank all the people who helped and supported me in writing this thesis.

Firstly, I would like to express my deepest thanks to my supervisor, Assoc. Prof. Dr. Syed Abdul Rahman Al-haddad B. Syed Mohamed, for his sincerity, patience and support. Undoubtedly, without his support, I never could finish my PhD journey. I am grateful for all the things he has done for me and thank God to give me the chance of being his student.

Secondly, I would like to appreciate the patience and supports of my co-supervisors: Prof. Dr. M. Iqbal Saripan and Prof. Dr. Shamala K Subramaniam. It was a great fortune of mine to attend in the courses lectured by them and be their student. They thought me a lot about the rules and disciplines of publishing papers and listened to my problems during my research friendly and with patience and for that I am very thankful.

Thirdly, I appreciate my husband for his understanding and cooperation. I know that he has tolerated many difficulties during my study while he always encouraged and guided me to accomplish my PhD degree.

Finally, I would like to appreciate supports of my friends in Multimedia Lab especially Nogol, Mehrdad and Zarina. I wish them success in all stages of their academic and family life.

This thesis was submitted to the Senate of Universiti Putra Malaysia and has been accepted as fulfilment of the requirement for the degree of Doctor of Philosophy. The members of the Supervisory Committee were as follows:

Syed Abdul Rahman Al-Haddad B. Syed Mohamed, PhD

Associate Professor
Faculty of Engineering
Universiti Putra Malaysia
(Chairman)

M. Iqbal Bin Saripan, PhD

Professor
Faculty of Engineering
Universiti Putra Malaysia
(Member)

Shamala Subramaniam, PhD

Professor
Faculty of Computer Science and Information Technology
Universiti Putra Malaysia
(Member)

BUJANG KIM HUAT, PhD

Professor and Dean
School of Graduate Studies
Universiti Putra Malaysia

Date:

Declaration by graduate student

I hereby confirm that:

- this thesis is my original work;
- quotations, illustrations and citations have been duly referenced;
- this thesis has not been submitted previously or concurrently for any other degree at any other institutions;
- intellectual property from the thesis and copyright of thesis are fully-owned by Universiti Putra Malaysia, as according to the Universiti Putra Malaysia (Research) Rules 2012;
- written permission must be obtained from supervisor and the office of Deputy Vice-Chancellor (Research and Innovation) before thesis is published (in the form of written, printed or in electronic form) including books, journals, modules, proceedings, popular writings, seminar papers, manuscripts, posters, reports, lecture notes, learning modules or any other materials as stated in the Universiti Putra Malaysia (Research) Rules 2012;
- there is no plagiarism or data falsification/fabrication in the thesis, and scholarly integrity is upheld as according to the Universiti Putra Malaysia (Graduate Studies) Rules 2003 (Revision 2012-2013) and the Universiti Putra Malaysia (Research) Rules 2012. The thesis has undergone plagiarism detection software.

Signature: _____ Date: _____

Name and Matric No.: _____

TABLE OF CONTENTS

	Page
ABSTRACT	i
ABSTRAK	iii
ACKNOWLEDGEMENTS	v
APPROVAL	vi
DECLARATION	viii
LIST OF TABLES	xii
LIST OF FIGURES	xiii
LIST OF ABBREVIATIONS	xv
 CHAPTER	
1 INTRODUCTION	1
1.1 Background	1
1.2 Problem Statement	2
1.3 Research Questions	4
1.4 Research Objectives	4
1.5 Contribution	5
1.6 Scope of Project	6
1.7 Research Outline	7
 2 LITERATURE REVIEW	9
2.1 Introduction	9
2.2 Distributed MutEx (DME) algorithm	10
2.2.1 Token-based algorithm	10
2.2.2 Permission-based algorithm	11
2.2.3 Hybrid algorithm	15
2.3 MANET MutEx Algorithms	16
2.3.1 Static MANET MutEx algorithm	16
2.3.2 Dynamic MANET MutEx algorithm	17
2.4 Conclusion	25
 3 METHODOLOGY	27
3.1 Introduction	27
3.2 Ancestral technique	29
3.3 System Description	31
3.3.1 Communication between nodes	31
3.3.2 Policy P	31
3.3.3 Failure of nodes or links	31
3.4 Performance Metrics	32
3.4.1 Availability	32
3.4.2 Messages per CS entry	32
3.4.3 Response time per CS entry	33
3.4.4 Level of Fault-tolerance	33
3.4.5 Energy efficiency	33
3.5 Experimental Setup	33

3.6	Distributed Ancestral Dynamic (DAD) Algorithm for Conventional DS	34
3.6.1	Data structure and timers	34
3.6.2	Initialization of Nodes	36
3.6.3	Event Handling	37
3.7	MANET Dynamic Ancestral Algorithm (MDA) for MANET	42
3.7.1	Data structure and timers	42
3.7.2	Initialization Procedure	47
3.7.3	Event Handler Procedure	47
3.8	Conclusion	52
4	RESULT AND DISCUSSION	53
4.1	Introduction	53
4.2	Performance Evaluation of Distributed Ancestral Dynamic (DAD) Algorithm	53
4.2.1	Proof of Correctness	53
4.2.2	Simulation Test Harness	55
4.2.3	Experimental Results	57
4.3	Performance evaluation of MANET Dynamic Ancestral Algorithm (MDA)	66
4.3.1	Proof of correctness	66
4.3.2	Simulation Test Harness	68
4.3.3	Experimental Results	70
4.4	Conclusion	77
5	CONCLUSION AND FUTURE WORK	79
5.1	Conclusions	79
5.2	Achievements	80
5.3	Limitations	81
5.4	Recommendation for Future Works	81
	REFERENCES	82
	APPENDICES	90
	BIODATA OF STUDENT	98
	LIST OF PUBLICATIONS	99

LIST OF TABLES

Table		Page
2-1	Summary of dynamic token-based MANET MutEx algorithms	21
2-2	Summary of dynamic non-token based MANET MutEx algorithms	25
3-1	Models supported by GloMoSim (Bajaj et al., 1997) and models used in simulation	29
4-1	Comparison of DAD, Dynamic Linear and Hybrid voting algorithms for 5 consecutive tests	59
4-2	Experimental setup configured in “config.in”	70
4-3	Simulation parameters for evaluation of MANET MutEx algorithms	71
4-4	The response time of PR algorithm (Wu, et al., 2008) and MDA algorithm for 4 and 8 nodes, high load and different levels of mobility	77

LIST OF FIGURES

Figure		Page
1-1	Relation between conventional distributed database and ad hoc database in a mobile database system (Opeoluwa & Opeyemi, 2010)	7
2-1	Simple classification of DME algorithms (Singh, et al., 2013)	10
2-2	Structure of basic Token-based algorithm through logical ring	11
2-3	Basic Permission-based algorithm (RA algorithm)	13
2-4	Classification of MANET MutEx algorithms	19
3-1	The overall methodology	28
3-2	Successful CS entry in MANET by using ancestral technique	30
3-3	An unsuccessful request for CS in MANET by using ancestral technique	30
3-4	Three procedures for development of a MutEx algorithm in DS and MANET	34
3-5	Relation between messages and node states in DAD algorithm	35
3-6	Relation between Request_timer and Release_timer	36
3-7	Procedure DAD_Initialize(Nk) sets initial local variables on every node Nk	37
3-8	Figure 3-8: Procedure DAD_CSing(Nk) generates CS requests and sends them to the nodes in partition Pk	38
3-9	Procedure DAD_EventHandler on every node	39
3-10	Procedure Update_request upon receiving Msg1 from the request establisher	40
3-11	Procedure Vote_reply, the establisher of CS request receives replies and decides based on incoming information	41
3-12	Procedure Commitment_release	42
3-13	States of a node regarding operation in CS and failure	43
3-14	Normal Execution of MDA algorithm	44
3-15	Execution of MDA algorithm when ancestor does not respond for a long time	45
3-16	Execution of MDA algorithm if ancestor waits for release beyond a finite time or a requester fails	46
3-17	Procedure MDA_Initialize (Nk) sets initial local variables on every node Nk	47
3-18	Sending REQUEST and REPLY in MANET	48
3-19	Procedure MDA_SendRequest(Nk) for sending the request	49
3-20	Procedure MDA_HandleRequest (Nj) for handling the incoming requests	50
3-21	Procedure MDA_HandleRelease (New_anc, tsanc) releases the CS and assigns new ancestor	50
3-22	Procedure MDA_HandleReply (Nk) replies to update requester	51
4-1	Simulation test harness used for MutEx algorithm in DS	56
4-2	A DS with two partitions	57
4-3	Availability comparison of DAD, Dynamic Linear and Hybrid voting algorithms	60

4-4	Availability of DAD and Hybrid voting algorithms for 50 run times	61
4-5	Mean availability of DAD, Dynamic Linear and Hybrid voting algorithms for 1000 runs	61
4-6	Availability of DAD, Dynamic Linear and Hybrid Voting algorithms relating to the number of failures, after 1000 runs for each scenario.	62
4-7	Availability of dynamic voting algorithms for low rate of failure in terms of number of nodes after 1000 runs for each scenario	62
4-8	Availability of dynamic voting algorithms for high rate of failure in terms of number of nodes, after 1000 runs for each scenario	63
4-9	Availability of voting algorithms for a partially connected DS consisting of 10 nodes after 1000 runs in respect to the number of failures	64
4-10	Availability of voting algorithms for a partially connected DS consisting of 100 nodes after 10000 runs in respect to the percentage of failures	64
4-11	Life time and availability of voting algorithms for consecutive failures. 10-100 nodes in DS and in each run, the maximum number of extra injected fault was 1. Each fault injection scenario was simulated 1000 times.	65
4-12	Time complexity of DAD versus Dynamic Linear and Hybrid algorithms for fully connected system in terms of percentages of failure	66
4-13	Simulation test harness for Mutex algorithm in MANET	69
4-14	MDA permission-based algorithm	71
4-15	Availability comparison of MDA and PR algorithm (Wu, et al., 2008) for different rate of mobility and low load	72
4-16	Availability comparison of MDA and PR algorithm (Wu, et al., 2008) for different mobility and high load	73
4-17	Availability of MDA algorithm for different rate of failure in terms of number of nodes for random rate of mobility	74
4-18	Response time of MDA algorithm for different rate of mobility and low load	76
4-19	Response time of MDA algorithm and PR algorithm (Wu, et al., 2008) for 4 nodes, different levels of mobility and low load	76
4-20	Response time of MDA algorithm for different levels of mobility and high load	77

LIST OF ABBREVIATIONS

Anc	Ancestor
AODV	Adhoc On-demand Distance Vector
BFA	Backbone Formation Algorithm
CS	Critical Section
CSMA	Carrier Sense Multiple Access
DAD	Distributed Ancestral Dynamic
DME	Distributed Mutual Exclusion
DN	Distinguished Nodes
DP	Distinguished Partition
DS	Distributed System
DSDV	Destination-Sequenced Distance Vector Routing
DSR	Dynamic Source Routing
ES	Entry Section
FIFO	First In First Out
FTP	File Transfer Protocol
GloMoSim	Global Mobile Information System Simulator
MANET	Mobile Ad hoc NETWORK
MCA	Merging Clustering Algorithm
MDA	MANET Dynamic Ancestral Algorithm
MH	Mobile Host
MSS	Mobile Support Station
MutEx	Mutual Exclusion
NC	Node Copies
NCS	Non-Critical Section
NS2	Network Simulator 2
PARSEC	Parallel Runtime Scheduling and Execution Controller
PR	Permission
RA	Ricart-Agrawala
RCV	Relative Consensus Voting
RL	Reverse Link
RSVP	Resource Reservation Protocol
TCP	Transmission Control Protocol
TORA	Temperory-Ordered Routing Algorithm
TTL	Time To Live
UDP	User Datagram Protocol
VN	Version Number
WiMAX	Worldwide Interoperability for Microwave Access

CHAPTER 1

INTRODUCTION

1.1 Background

A Mobile Ad Hoc Network (MANET) is a collection of autonomous wireless mobile nodes established on a non-fixed infrastructure (Fife & Gruenwald, 2003). MANET is a kind of distributed system (DS) (Challenger et al., 2013) where links are wireless and each mobile node can communicate with its single or multi-hop neighbors through message passing. Message passing is the sole means for implementing and timestamp (Lamport, 1978a; Ricart & Agrawala, 1981) is basic strategy for prioritizing the requests for critical sections in DS and its extensions, e.g. mobile ad hoc networks. Due to absence of a central coordinator, each mobile node acts as a router.

Due to mobility of nodes and frequent changes in topology of network, mobile nodes and wireless links are vulnerable to fail. Failure of nodes or links yields to transient or permanent failure of Mutex algorithms. Fault-tolerant Mutex algorithms utilize different methods to deal with failure of nodes and links, while some algorithms try to detect and repair faulty nodes, other group are equipped with techniques, e.g., redundancy, to tolerate faults. Redundancy has been widely utilized to improve fault-tolerance of hardware or software systems in four main form of redundancy of software version, hardware modules, time and data. Replication of data improves the fault-tolerance of database systems. MANET and other kinds of DS use the replication of one or more data resources in the following ways: to increase the system's availability, to service requests for the same information in parallel (Gifford, 1979) to decrease the system's response time and communication costs by providing the nearest copy of the resource to the location of demand, and to have load sharing by distributing the computational load of responding to queries among a number of nodes rather than centralized in a single node (Thomas, 1979). Besides these advantages, there is a huge challenge in the maintenance of the replicated resources in such a way that at the most, the one node can access the shared resource(s) at a given time t . This problem is well known as Mutual Exclusion (Mutex), and is still an active research area since more than three decades ago when the first distributed Mutex (DME) algorithm presentations established by (Le Lann, 1977) and (Lamport, 1978b). In order to facilitate efficient data access and update, databases are deployed on MANETs (Padmanabhan et al., 2008). Data replication is widely used in MANET database (Opeoluwa & Opeyemi, 2010) which improves the availability of data (Lipskoch & Theel, 2014).

Many reasons make the conventional Mutex algorithms inefficient and even infeasible in MANET. Due to mobility of nodes, the geographic position and topology of MANET change very frequently. Although, knowing the topology of network leads to improvements in Mutex algorithms in terms of the number of messages transmitted, energy saving, time complexity of algorithms and the order of visiting the nodes, it is very expensive for MANETs (Malpani et al., 2005). Due to mobility and wireless communication of nodes, crash of software or hardware, node movement to out of radius range, power and battery constraints, and no fixed client/server or coordinator, DME algorithms cannot be simply implemented on

MANET. Because of mobility and very frequent changes in the topology, the method of formation or removal of wireless links is unpredictable. Mobile nodes are also limited by the energy and processing power compared to fixed nodes. A higher mobility rate and more radio transmission increase the power consumption of mobile nodes, which leads to a lack of battery and the failure of nodes. Furthermore, due to a higher probability of node or link failures (Moallemi et al., 2007) and a high ratio of faults in wireless connections, MANETs are relatively high in failure rate compared to wired linked networks. Mobile Support Station (MSS) in conventional DS can act on behalf of the mobile hosts (MH) once failure occurs; however, there are no infrastructure and MSS in MANETs (Wu et al., 2005).

Conventional algorithms for MutEx in DS are categorized into two main classes including token-based and permission-based algorithms. This classification is presented for MANET as well. Token-based algorithms depend on circulation of a specific message known as token. The owner of the token has priority for entering the Critical Section (CS). They required logical structure (tree or ring), and useless circulation of token when there is no request for CS. Token may be lost during communications, because of link failure or failure of token host. Unfortunately, the process of recovering a new token after token loss is very complicated and time consuming in MANETs in comparison with conventional DS. Because of common problem of single token as a single point of failure, token-based algorithms are generally non-fault-tolerant (although some mechanisms are utilized to increase their level of fault-tolerance). Almost all MANET MutEx algorithms are token-based (except (Masum et al., 2010; Wu et al., 2008)). Token-based algorithms (Attiya et al., 2010; Baldoni et al., 2002; Jiang, 2003; Moallemi, et al., 2007; Tamhane & Kumar, 2011; Yang, 2005) have some beneficial features for MANETs including fewer number of messages transmitted and less information to be stored about other mobile nodes.

Permission-based algorithms utilize the permission of multiple nodes to guarantee MutEx. Basic form of permission-based MutEx algorithm permits a CS request to enter the CS when permission of at least the majority of nodes is received by requesting node. Permission-based algorithms have some unique features including there is no need to maintain logical topology to pass the token, neither to propagate any message if no host requests to enter critical section (Wu, et al., 2008). However, it leads to high number of message transmissions and increments energy consumption in MANET by increasing the number of mobile nodes.

1.2 Problem Statement

MutEx is a well-known problem in maintenance of replicated data in DS, in which at a given time t at most one node is eligible to update common data (Iakab, 2012). MutEx becomes more complicated in MANETs, because topology and position of mobile nodes changes very frequently (Erciyes & Dagdeviren, 2012), wireless links are more vulnerable to fail, there is no central coordinator, and every node is responsible for routing the messages (Gupta et al., 2012).

This study deals with the following research problems:

- a) Consecutive failure of mobile nodes or links causes loss of a large number of small packets and a high traffic on routing protocols (Gill et al., 2011) which direct the messages from upper layer of network (here application layer) to destination nodes through communication channels. Obviously, the traffic on routing layer is caused by many factors; however, the factor related to this study is the number of messages generated by MutEx algorithm. It is important to decrease the number of message transmissions per CS entry as introduced in literature. In order to measure the number of messages per CS entry, the number of message transmissions is counted by a node since it has issued the request for the CS until it receives the CS permission (or token).
- b) For every unanswered request for CS, reattempts are performed by requesting nodes (Parameswaran & Hota, 2010). Every packet has a limited Time To Live (TTL) and also a limited and predefined time-out is considered for every request for CS (Mallikarjuna et al., 2012). Due to time-outs in the nodes, failure of nodes or links leads to unanswered (or unsuccessful) CS request. Obviously, with higher successful CS entries, lower reattempts are performed and processes by nodes. Service availability (Gill et al., 2011) (henceforth, name as availability) is simply defined as the probability that an update request can be serviced at any time, t . The ratio of successful CS entries over network lifetime simply reflects the probability that an update request can be serviced successfully by DS or simply the availability of MutEx. One solution for high reattempts is to degradation in the number of unsuccessful CS entries, i.e., improving the rate of successful CS entries or simply improving the availability of CS.
- c) Response time is the time interval for a node to issue its request, until it enters the CS (Sharma et al., 2014; Wu, et al., 2008). Permission-based algorithm generally wait until they receive permission from a predefined set or number of nodes. In case of failure of one or more nodes, the node is halted until the faulty node is repaired. This event increases the response time to CS request or response time per CS entry. Therefore, it is important to find an approach to decrease response time per CS entry.
- d) Partitioning is a kind of failure that may occur in the network due to failure of some mobile hosts or links (Shi & Chen, 2014). When a network is partitioned, it is divided to two or more isolated groups of nodes that can communicate inside the groups but they cannot communicate with other nodes in other groups. Hence, each partition does not have any view from the other partition. Partitioning is harmful for MutEx if two nodes in two partitions decide to enter the CS simultaneously (Wu et al., 2008). In order to face with partitioning, MutEx algorithm should have a strategy to prevent the entrance of isolated partitions to critical section at the same time. This problem has been studied in DS and is a major issue that should especially be considered in MANET (Shi & Chen, 2014). Due to failure of mobile nodes or wireless links, disconnection of nodes, movement of nodes to out of radius range, or battery deployment of nodes, MANET is more likely to be partitioned to isolated groups of nodes. Thus, a MANET MutEx algorithm,

as well as conventional DS Mutex, should be able to tolerate faults and transient partitioning of network.

1.3 Research Questions

This research embarks on following questions:

- a) Can be found any solution for Mutex, other than permission and token, to resolve problem of high traffic in permission-based algorithms and problem of token as a single point of failure?
- b) How new solution improves permission-based Mutex?
- c) How the problem of token as a single point of failure is resolved?
- d) How is the scalability of new solution when number of nodes becomes large?
- e) Do MANET and DS tolerate failure of nodes or links?
- f) How long a MANET and DS can resist against failure of nodes or links?
- g) How is its fault-tolerant for different scenarios of failure occurrence?
- h) How is the performance of new solution in terms of availability, time complexity and message complexity?

1.4 Research Objectives

The main purpose of this research is to introduce a method of managing the CS having fault-tolerant and partition-tolerant features with fewer message transmissions and response time than other permission-based algorithms. The research objectives are as follows:

- a) To design a permission-based technique to enhance the fault tolerance of token-based and resolve the high number of message transmissions of previous permission-based algorithms.
- b) To develop a fault-tolerant and permission-based algorithm to resolve the problem of Mutex in conventional DS, which prolongs the service to CS request and improves the availability of CS.
- c) To develop a fault-tolerant and permission-based algorithm for resolving the problem of Mutex in MANET with higher availability of CS, lower response time for a CS request to be served and fewer numbers of messages imposed to network.

1.5 Contribution

The topic of Mutex algorithms for MANET is an ongoing topic. In this study, a review on dynamic Mutex algorithms is presented for MANETs along with a classification on the techniques involved in decision making. It is hoped that this discussion can increase the understanding on this subject and help the designers of systems and engineers to implement efficient protocols for MANETs. This study reviews Mutex algorithms in different way with (Benchaiiba et al., 2004).

Ancestral method as used in this research resolves the issues related to token-based and permission-based algorithms. This technique has been first introduced in this study. Two algorithms based on ancestral technique are introduced and evaluated in this study. The first algorithm (DAD algorithm) resolves the problem of Mutex in DS. It also improves the availability of CS and prolongs the life time of service. Second Ancestral Mutex algorithm (MDA algorithm) is presented for MANET and evaluated for different levels of mobility of nodes, fault-injection and number of nodes. It guarantees Mutex, dead lock freedom and starvation freedom and can be used in every situation where mobile devices are connected each other in form of MANET and helps them to stay up when multiple failure of mobile nodes or links occurs. It resolves the problem of network partitioning which frequently occurs in DS and MANET.

Ancestral technique is different with combination of token and permission based technique which has been introduced as hybrid technique. It does not utilize any token, that is similar to permission-based, and the latest node having the CS influences the entrance of the next node to the CS, that is similar to token-based algorithms. As the technique needs permission of other node, i.e., it can be counted as permission-based. Fault tolerance or even partition tolerance feature of ancestral method, keeps MANET up in systems such as disaster management systems (Gupta, et al, 2012) (for discovering alive human or their body) or traffic control system (Luo et al., 2014) (for avoiding heavy traffic on roads or tunnels).

By decreasing the response time per CS entry and increasing the availability of CS, MDA algorithm is more efficient than similar methods. It needs fewer messages to be transmitted between mobile devices in order to access CS. This feature influences the energy efficiency of MANET because higher energy will be preserved and saved during wireless communication of mobile devices. MDA algorithm can be utilized where the mobile nodes need to share resources or codes of the data in all of these applications. Mutex is required to guarantee the consistency of the shared objects in automated battlefields, forestry, environmental analysis, disaster managements and traffic control applications.

MDA algorithm does not need to know every connected node and works according to local knowledge of a node from its neighbors. Therefore, MDA is more realistic. MDA algorithm benefits the main assumption of previous permission-based algorithms on MANET environment; however, they are basically different.

1.6 Scope of the Project

The project focuses on fault-tolerant MutEx algorithms for DS and MANET. Among different categories of fault-tolerant MutEx algorithms, permission-based algorithms are concentrated for resolving the problem of MutEx in conventional DS and MANET as a kind of DS (Challenger et al., 2013). Nowadays, mobile systems are used everywhere. A decade ago, almost all database systems were established in fixed location in different geographic positions. Fixed database systems are not efficient anymore regarding the widespread using of mobile devices, cellular phones, and Laptops. Furthermore, WiFi and WiMax provide internet connection almost everywhere. Therefore, using mobile Ad Hoc database systems is feasible in many situations. A direct application of presented algorithm in this study is in mobile Ad Hoc database systems.

A mobile Ad Hoc database may be used with or without a conventional database system. In Figure 1-1, a system with both conventional and Ad Hoc database system is illustrated. In order to facilitate efficient data access and update, databases are deployed on MANETs (Azeem & Khan, 2012; Padmanabhan, et al., 2008). An ideal database system must provide the most up-to-date data to every request for inquiry which is only possible when replicated data is up-to-date. MutEx guarantees that at a given time t only one node updates shared data, while data consistency guarantees that every node in network updates its copy from shared data to the latest version. This study is limited to MutEx algorithm and assumptions related to data consistency are not taken into account. Mobility of nodes makes MANET sustainable to multiple connection or disconnection of nodes, in addition to failure of nodes. MANET as concentrated in this study is fault-prone and its topology changes frequently.

The DS and the MANET as concentrated in this study initially form one partition P ; however, some nodes may be timely disconnected from the partition due to failures. Every single node N_k sees the partition from its point of view as partition P_k , so that initially $P_k = P$. In the case of partitioning, some partitions may have outdated data (because they do not receive updates from the DP), and reading from such partitions may result in outdated information. Nodes do not need to know the information of other partitions and only act based on their partition information. However, concurrent read operations are allowed if the node belongs to a DP. Therefore, for each query, reading permission from one node is enough (Barbara & Garcia-Molina, 1986). This study concentrated on writing on replicated data or update.

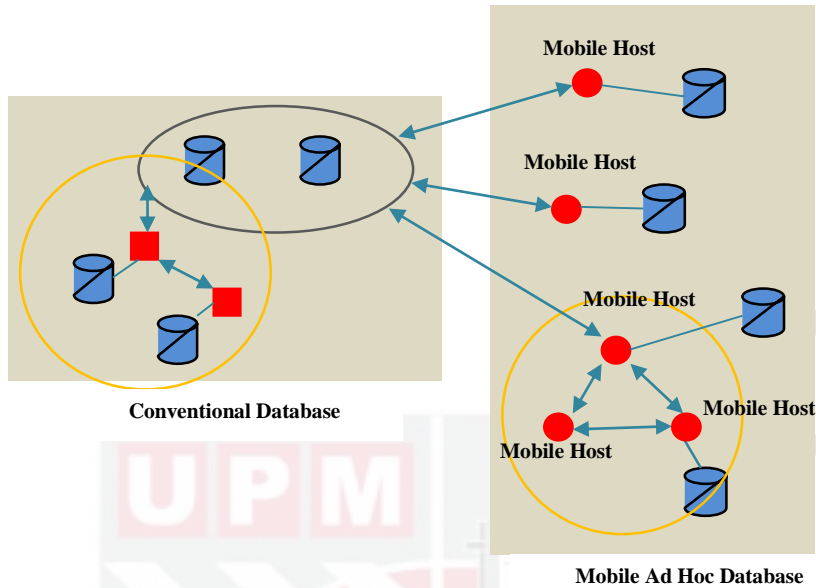


Figure 1-1: Relation between conventional distributed database and ad hoc database in a mobile database system (Opeoluwa & Opeyemi, 2010)

1.7 Research Outline

Chapter 1 introduces the problem of MutEx in DS (and database system) and its relation to mobile data base system as used in mobile ad hoc environments. Two main classes of MutEx algorithms for DS and MANET as a generalization of conventional DS are explained and their benefits and problems are highlighted. The research objectives are stated and the problem statement in current situation and contribution are also given.

Chapter 2 is dedicated to literature review on fault-tolerant MutEx algorithms and begins with introduction to MutEx algorithms. DME algorithms related to MANET algorithms are discussed. The benefits and weaknesses of each algorithm are also highlighted. Finally, MutEx DS and MANET algorithms are summarized and compared.

Chapter 3 describes the system model, performance metrics and methodology used for designation of DAD and MDA algorithms. It also presents the description of basic distributed algorithm and the MANET MutEx algorithms. The experimental setup that is used for development and evaluation of the algorithms is also presented in this chapter.

Chapter 4 discusses the result of distributed algorithms in conventional DS and in MANET. The MutEx, deadlock freedom and starvation freedom of algorithms are proved. Different scenarios of fault injection, different speed of mobility of nodes and load of requests are considered for the purpose of simulation.

Chapter 5 concludes the study by summarizing the result and conclusion, stating the contributions and identifying future research direction.



REFERENCES

- Agrawal, D., & Abbadi, A. E. (1991). An Efficient and Fault-tolerant Solution for Distributed Mutual Exclusion. *ACM Transactions on Computing Systems*, 9(1), 1-20.
- Agrawal, G., & Jalote, P. (1992). *An Efficient Protocol for Voting in Distributed Systems*. Paper presented at the 12th International Conference on Distributed Computing Systems.
- Alsberg, P. A., & Day, J. D. (1976). *A principle for resilient sharing of distributed resources*. Paper presented at the Proceedings of the 2nd international conference on Software engineering, San Francisco, California, United States.
- Anchal, Poonam Saini, & Krishna, C. R. (2014). *An Efficient Permission-Cum-Cluster Based Distributed Mutual Exclusion Algorithm for Mobile Adhoc Networks*. Paper presented at the 2014 International Conference on Parallel, Distributed and Grid Computing.
- Attiya, H., Kogan, A., & Welch, J. L. (2008). *Efficient and Robust Local Mutual Exclusion in Mobile Ad Hoc Networks*. Paper presented at the The 28th International Conference on Distributed Computing Systems, 2008. ICDCS '08. .
- Attiya, H., Kogan, A., & Welch, J. L. (2010). Efficient and Robust Local Mutual Exclusion in Mobile Ad Hoc Networks. *IEEE Transactions on Mobile Computing*, 9(3), 361-375.
- Azeem, R., & Khan, M. I. A. (2012). Techniques about Data Replication for Mobile Ad-hoc Network Databases. *International Journal Of Multidisciplinary Sciences And Engineering*, 3(5), 53-57.
- Bajaj, L., Takai, M., Ahuja, R., Tang, K., Bagrodia, R., & Gerla, M. (1997). GloMoSim: A scalable network simulation environment: Univ. California
- Baldoni, R., Virgillito, A., & Petrassi, R. (2002). *A distributed mutual exclusion algorithm for mobile ad-hoc networks*. Paper presented at the Proceedings of the Seventh International Symposium on Computers and Communications (ISCC 2002), 2002. .
- Barbara, D., & Garcia-Molina, H. (1986). Mutual exclusion in partitioned distributed systems. *Distributed Computing*, 1(2), 119-132.
- Benchaiiba, M., Bouabdallah, A., Badache, N., & Nacer, M. A. (2004). Distributed mutual exclusion algorithms in mobile ad hoc networks: an overview. *SIGOPS Operating Systems Review*, 38(1), 74-89.

- Benchaiiba, M., & Nacer, M. A. (2010). A distributed token based h-out of-k Mutual Exclusion protocol for mobile ad hoc networks. *International Journal of Ad Hoc Ubiquitous Computing*, 5(2), 117-135.
- Birman, K. P. (2012). *Guide to reliable distributed systems* (Vol. XXII): Springer.
- Budhiraja, N., Marzullo, K., Schneider, F. B., & Toueg, S. (1993). *The primary-backup approach Distributed systems (2nd ed.)*: ACM Press/Addison-Wesley Publishing Co.
- Challenger, M., Haytaoglu, E., Tokatli, G., Dagdeviren, O., & Erciyes, K. (2013). A Hybrid Distributed Mutual Exclusion Algorithm for Cluster-Based Systems. *Mathematical Problems in Engineering*, 2013, 15.
- Chang, Y. I., Singhal, M., & Liu, M. T. (1990). *A fault tolerant algorithm for distributed mutual exclusion*. Paper presented at the Proceedings of the Ninth Symposium on Reliable Distributed Systems, 1990. .
- Cheung, S. Y., Ahamad, M., & Ammar, M. H. (1989). Optimizing vote and quorum assignments for reading and writing replicated data. *IEEE Transactions on Knowledge and Data Engineering*, 1(3), 387-397.
- Coulouris, G., Dollimore, J., & Kindberg, T. (2005). *Distributed Systems: Concepts and Design* (fourth ed.): Addison-Wesley.
- Dagdeviren, O., & Erciyes, K. (2006). *A distributed backbone formation algorithm for mobile ad hoc networks*. Paper presented at the Proceedings of the 4th international conference on Parallel and Distributed Processing and Applications, Sorrento, Italy.
- Dagdeviren, O., Erciyes, K., & Cokuslu, D. (2005). *Merging clustering algorithms in mobile ad hoc networks*. Paper presented at the Proceedings of the Second international conference on Distributed Computing and Internet Technology, Bhubaneswar, India.
- Davcev, D. (1989). A Dynamic Voting Scheme in Distributed Systems. *IEEE Transactions on Software Engineering*, 15(1), 93-97.
- Derhab, A., & Zair, M. (2010). *A resource-based mutual exclusion algorithm supporting dynamic acting range and mobility for Wireless sensor and Actor Networks*. Paper presented at the 6th IEEE International Conference on Distributed Computing in Sensor Systems Workshops (DCOSSW), 2010
- Dhamdhare, D. M., & Kulkarni, S. S. (1994). A token based k-resilient mutual exclusion algorithm for distributed systems. *Information Processing Letters*, 50(3), 151-157.
- Dijkstra, E. W. (1965). Co-operating Sequential Processes.

- Erciyes, K. (2004). *Cluster based Distributed Mutual Exclusion Algorithms for Mobile Networks*. Paper presented at the Euro-Par 2004.
- Erciyes, K., & Dagdeviren, O. (2012). A Distributed Mutual Exclusion Algorithm for Mobile Ad Hoc Networks. *International Journal of Computer Networks & Communications*, 4(2), 129-148.
- Fife, L. D., & Gruenwald, L. (2003). Research issues for data communication in mobile ad-hoc network database systems. *SIGMOD Rec.*, 32(2), 42-47.
- Fischer, M. J., Lynch, N. A., & Paterson, M. S. (1985). Impossibility of Distributed Consensus with One Faulty Process. *Journal of the ACM*, 32(2), 374-382.
- Gafni, E., & Bertsekas, D. (1981). Distributed Algorithms for Generating Loop-Free Routes in Networks with Frequently Changing Topology. *IEEE Transactions on Communications*, 29(1), 11-18.
- Garcia-Molina, H., & Barbara, D. (1985). How to assign votes in a distributed system. *Journal of the ACM*, 32(4), 841-860.
- Gifford, D. K. (1979). *Weighted Voting for Replicated Data*. Paper presented at the Proceedings of the seventh ACM symposium on Operating systems principles, Pacific Grove, California, United States.
- Gill, P., Jain, N., & Nagappan, N. (2011). Understanding network failures in data centers: measurement, analysis, and implications. *SIGCOMM Comput. Commun. Rev.*, 41(4), 350-361.
- Gupta, A., Reddy, B. V. R., Ghosh, U., & Khanna, A. (2012). A Permission-based Clustering Mutual Exclusion Algorithm for Mobile Ad-Hoc Networks. *International Journal of Engineering Research and Applications*, 2(4), 019-026.
- Hoare, C. A. R. (1974). Monitors: an operating system structuring concept. *Commun. ACM*, 17(10), 549-557.
- Iakab, K. K. (2012). *Probabilistic Quorum Systems for Dependable Distributed Data Management*. Carl Von Ossietzky universitat.
- Ingols, K., & Keidar, I. (2001). *Availability study of dynamic voting algorithms*. Paper presented at the 21st International Conference on Distributed Computing Systems.
- Jajodia, S., & Mutchler, D. (1987). Dynamic Voting. *SIGMOD Rec.*, 16(3), 227-238.
- Jajodia, S., & Mutchler, D. (1989). A Hybrid Replica Control Algorithm Combining Static and Dynamic Voting. *IEEE Transactions on Knowledge and Data Engineering*, 1(4), 459-469.

- Jajodia, S., & Mutchler, D. (1990). Dynamic Voting Algorithms for Maintaining the Consistency of a Replicated Database. *ACM Transactions on Database Systems*, 15(2), 230-280.
- Jiang, J.-R. (2003). *A prioritized h-out-of-k mutual exclusion algorithm with maximum degree of concurrency for mobile ad hoc networks and distributed systems*. Paper presented at the Proceedings of the Fourth International Conference on Parallel and Distributed Computing, Applications and Technologies, 2003. PDCAT'2003. .
- Jiannong, C., Jingyang, Z., Daoxu, C., & Wu, J. (2004). *An Efficient Distributed Mutual Exclusion Algorithm based on Relative Consensus Voting*. Paper presented at the Proceedings of 18th International Parallel and Distributed Processing Symposium
- Kakugawa, H., Fujita, S., Yamashita, M., & Ae, T. (1994). A distributed k-mutual exclusion algorithm using k-coterie. *Information Processing Letters*, 49(4), 213-218.
- Karthick, R., & Gopinathan, B. (2011). Efficient Algorithms for Distributed Mutual Exclusion in Mobile Ad-Hoc Network. *Indian Journal of Computer Science and Engineering*, 2(6), 873-884.
- Khanvilkar, S., & Shatz, S. M. (2001). Tool integration for flexible simulation of distributed algorithms. *Software—Practice and experience*, 31, 1363–1380.
- Kumar, A. (1991). Hierarchical Quorum Consensus: a New Algorithm for Managing Replicated Data. *IEEE Transactions on Computers*, 40(9), 996-1004.
- Lamport, L. (1978a). The implementation of reliable distributed multiprocess systems. *Computer Networks (1976)*, 2(2), 95-114.
- Lamport, L. (1978b). Time, Clocks and the ordering of events in a distributed system. *Communications of the ACM*, 21(7), 558-565.
- Latif-Shabgahi, G., Bennett, S., & Bass, J. M. (2003). Smoothing voter: a novel voting algorithm for handling multiple errors in fault-tolerant control systems *Microprocessors and Microsystems*, 27(7), 303-313.
- Le Lann, G. (1977). *Distributed systems: towards of a formal approach*. Paper presented at the IFIP Congress, North-Holland.
- Lipskoch, K., & Theel, O. (2014). Relaxing data consistency along different dimensions for increasing operation availabilities. *International Journal of Parallel, Emergent and Distributed Systems*, 1-29.

- Luo, A., Cao, J., & Zhang, J. (2014). Distributed Mutual Exclusion Algorithms for Intersection Traffic Control. *IEEE Transactions on Parallel and Distributed Systems*, 1.
- Maekawa, M. (1985). A Sqrt(N) Algorithm for Mutual Exclusion in Decentralized Systems. *ACM Transactions on Computer Systems*, 3(2), 145-159.
- Mallikarjuna, B., Pusalatha, C., & Saritha, V. (2012). Fault Tolerant Resource Management with Mutual Exclusion Algorithm for Mobile Adhoc Networks. *International Journal of Scientific and Research Publications*, 2(4).
- Malpani, N., Vaidya, N. H., & Welch, J. L. (2001). *Distributed token circulation on mobile ad hoc networks*. Paper presented at the Ninth International Conference on Network Protocols, 2001.
- Malpani, N., Yu, C., Vaidya, N. H., & Welch, J. L. (2005). Distributed token circulation in mobile ad hoc networks. *IEEE Transactions on Mobile Computing*, 4(2), 154-165.
- Masum, S. M., Akbar, M. M., Ali, A. A., & Rahman, M. A. (2010). A consensus-based ℓ -Exclusion algorithm for mobile ad hoc networks. *Ad Hoc Networks*, 8(1), 30-45.
- Mellier, R., & Myoupo, J.-F. (2005). *A clustering mutual exclusion protocol for multi-hop mobile ad hoc networks*. Paper presented at the 13th IEEE International Conference on Networks, 2005. Jointly held with the 2005 IEEE 7th Malaysia International Conference on Communication, 2005.
- Mellier, R., Myoupo, J.-F., & Ravelomanana, V. (2005). A Non-Token-Based-Distributed Mutual Exclusion Algorithm for Single-Hop Mobile Ad Hoc Networks. In E. Belding-Royer, K. Al Agha & G. Pujolle (Eds.), *Mobile and Wireless Communication Networks* (Vol. 162, pp. 287-298): Springer US.
- Moallemi, M., Moghaddam, M. H. Y., & Naghibzadeh, M. (2007). *A Fault-Tolerant Mutual Exclusion Resource Reservation Protocol for Clustered Mobile Ad hoc Networks*. Paper presented at the Eighth ACIS International Conference on Software Engineering, Artificial Intelligence, Networking, and Parallel/Distributed Computing, 2007. SNPD 2007.
- Naimi, M., & Trehel, M. (1987). *A Distributed Algorithm for Mutual Exclusion based on Data Structures and Fault Tolerance*. Paper presented at the the 6th international Phoenix IEEE conference on computer and communication, Scottsdale.
- Opeoluwa, O. A., & Opeyemi, O. G. (2010). From conventional to mobile database management system: A theoretical review. *Journal of Mobile Communication*, 4(4), 96-99.

- Osrael, J., Froihofer, L., Chlaupke, N., & Goeschka, K. M. (2007). *Availability and Performance of the Adaptive Voting Replication*. Paper presented at the The Second International Conference on Availability, Reliability and Security, 2007. ARES 2007. .
- Osrael, J., Froihofer, L., & Goeschka, K. M. (2007). *Availability/Consistency Balancing Replication Model*. Paper presented at the IEEE International Parallel and Distributed Processing Symposium, 2007. IPDPS 2007. .
- Padmanabhan, P., Gruenwald, L., Vallur, A., & Atiquzzaman, M. (2008). A survey of data replication techniques for mobile ad hoc network databases. *The VLDB Journal*, 17(5), 1143-1164.
- Parameswaran, M., & Hota, C. (2010). *A novel permission-based reliable distributed mutual exclusion algorithm for manets*. . Paper presented at the 2010 Seventh International Conference On Wireless and Optical Communications Networks (WOCN), Hilton Colombo Hotel, Colombo, Sri Lanka.
- Park, S. H., Kim, Y., Han, J. H., & Park, J. S. (2010). *Quorum based mutual exclusion in asynchronous systems with unreliable failure detectors*. Paper presented at the In ITNG2010 - 7th International Conference on Information Technology: New Generations.
- Park, S. H., & Lee, S. H. (2014). Quorum-based mutual exclusion in asynchronous distributed systems with unreliable failure detectors. *Journal of Supercomputing*, 67, 469–484.
- Park, V. D., & Corson, M. S. (1997). *A Highly Adaptive Distributed Routing Algorithm for Mobile Wireless Networks*. Paper presented at the Proceedings of the INFOCOM '97. Sixteenth Annual Joint Conference of the IEEE Computer and Communications Societies. Driving the Information Revolution.
- Paydar, S., Naghibzadeh, M., & Yavari, A. (2006). *A Hybrid Distributed Mutual Exclusion Algorithm*. Paper presented at the ICET '06. International Conference on Emerging Technologies.
- Rahman, M. A., & Akbar, M. M. (2010). A Permission Based Hierarchical Algorithm for Mutual Exclusion. *Journal of Computers*, 5(12).
- Ravindran, K., Kwiat, K. A., & Ding, G. (2005). *Simulation-Based Validation of Protocols for Distributed Systems*. Paper presented at the System Sciences, 2005. HICSS '05. Proceedings of the 38th Annual Hawaii International Conference on.
- Raymond, K. (1989). A tree based algorithm for distributed mutual exclusion algorithms. *ACM Transactions on Computer Systems*, 7(1), 61- 77.

- Raynal, M. (1991). A simple taxonomy for distributed mutual exclusion algorithms. *SIGOPS Operating Systems Review*, 25(2), 47-50.
- Ricart, G., & Agrawala, A. K. (1981). An Optimal Algorithm for Mutual Exclusion in Computer Networks. *Communications*, 24(1), 9-17.
- Robinson, K., Turgut, D., & Chatterjee, M. (2006). *An Entropy-based Clustering in Mobile Ad hoc Networks*. Paper presented at the Proceedings of the 2006 IEEE International Conference on Networking, Sensing and Control, 2006. ICNSC '06.
- Sawai, Y., Shinohara, M., Kanzaki, A., Hara, T., & Nishio, S. (2007). *Quorum-based consistency management among replicas in ad hoc networks with data update*. Paper presented at the Proceedings of the 2007 ACM symposium on Applied computing, Seoul, Korea.
- Saxena, P. C., & Rai, J. (2003). A Survey of Permission-based Distributed Mutual Exclusion Algorithms. *Computer Standards & Interfaces*, 25(2), 159-181.
- Sharma, B., Bhatia, R. S., & Singh, A. K. (2014). A Token Based Protocol for Mutual Exclusion in Mobile Ad Hoc Networks. *J Inf Process Syst*, 10(1), 36-54.
- Shi, K., & Chen, H. (2014). RHPMAN: Replication in Highly Partitioned Mobile Ad Hoc Networks. *International Journal of Distributed Sensor Networks*, 2014, 14.
- Sil, S., & Das, S. (2010). An energy efficient Algorithm for Distributed Mutual Exclusion in Mobile Ad-hoc Networks. *World Academy of Science, Engineering and Technology*, 64, 517-522.
- Singh, R., Malviya, J., & Jha, K. (2013). A survey of mutual exclusion algorithms in distributed computing. *International Journal on Emerging Technologies*, 4(2), 124-127.
- Singhal, M. (1993). A Taxonomy of Distributed Mutual Exclusion *Journal of Parallel and Distributed Computing*, 18, 94-101.
- Singhal, M., & Manivannan, D. (1997). *A distributed mutual exclusion algorithm for mobile computing environments*. Paper presented at the Proceedings Intelligent Information Systems, 1997. IIS '97.
- Sung-Hoon, P., Tae-Gyu, L., Hyung-Seok, S., Seok-Jin, K., & Jong-Ho, H. (2009). *An Election Protocol in Mobile Ad Hoc Distributed Systems*. Paper presented at the Sixth International Conference on Information Technology: New Generations, 2009. ITNG '09. .
- Suzuki, I., & Kasami, T. (1985). A distributed mutual exclusion algorithm. *ACM Transactions on Computer Systems*, 3(4), 344-349.

- Swaroop, A., & Singh, A. K. (2007). A study of token based algorithms for distributed mutual exclusion *International Review on Computers & Software*, 2(4), 302.
- Tamhane, S. A., & Kumar, M. (2011). A token based distributed algorithm for supporting mutual exclusion in opportunistic networks. *Pervasive and Mobile Computing*(0).
- Thomas, R. H. (1979). A Majority Consensus Approach to Concurrency Control for Multiple Copy Databases. *ACM Transactions on Database Systems*, 4(2), 180-209.
- Tong, Z., & Kain, R. Y. (1991). Vote Assignments in Weighted Voting Mechanisms. *IEEE Transactions on Computers*, 40(5), 664-667.
- Vedantham, R., Zhenyun, Z., & Sivakumar, R. (2006). *Mutual Exclusion in Wireless Sensor and Actor Networks*. Paper presented at the 3rd Annual IEEE Communications Society on Sensor and Ad Hoc Communications and Networks, 2006. SECON '06.
- Velazquez, M. G. (1993). A survey of distributed mutual exclusion algorithms (Vol. Technical Report CS-93-116).
- Walter, J. E., & Kini, S. (1997). *Mutual Exclusion on Multihop, Mobile Wireless Networks*. Texas: Texas A & M University.
- Walter, J. E., Welch, J. L., & Vaidya, N. H. (2001). A mutual exclusion algorithm for ad hoc mobile networks. *Wireless Networks*, 7(6), 585-600.
- Wu, W., Cao, J., & Raynal, M. (2007). *A dual-token-based fault tolerant mutual exclusion algorithm for MANETs*. Paper presented at the Proceedings of the 3rd international conference on Mobile ad-hoc and sensor networks, Beijing, China.
- Wu, W., Cao, J., & Yang, J. (2005). *A scalable mutual exclusion algorithm for mobile ad hoc networks*. Paper presented at the Proceedings of the 14th International Conference on Computer Communications and Networks, ICCCN 2005.
- Wu, W., Cao, J., & Yang, J. (2008). A fault tolerant mutual exclusion algorithm for mobile ad hoc networks. *Pervasive and Mobile Computing*, 4(1), 139-160.
- Yang, C.-Z. (2005). *A token-based h-out of-k distributed mutual exclusion algorithm for mobile ad hoc networks*. Paper presented at the 3rd International Conference on Information Technology: Research and Education. ITRE 2005.
- Zhu, Y., Zhang, H., & Ji, Q. (2013). How much delay has to be tolerated in a mobile social network? *International Journal of Distributed Sensor Networks*, 2013(